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#### DATA TRANSMITTING METHOD (54)

An art for minimizing transmission delay of lowrate data such as voice data and for enhancing the transmission efficiency in multiplexed data communication using with AAL type 2 is disclosed. The SSCS of the high-rate data side of the AAL type 2 is responsible for addition and elimination of a dummy code field, error correction, and division and reassembly of data unit. Therefor, a dummy code field 3 and restorative code field 10 are attached to a user data unit 1 which is supplied from the higher layer in a source device, whereby a data unit 11 is prepared. The data unit 11 is then divided into segments 12-1 to 12-m of a fixed-length, 45 octets. Successive code fields 6-1 to 6-m are attached to the segments 12-1 to 12-m, whereby data blocks 13-1 to 13-m are prepared. The data blocks 13-1 to 13-m are successively transferred to an AAL-2 converter 24. In a destination device, from the data blocks 13-1 to 13m, the user data unit 1 is reassembled conversely.

FIG. 1A

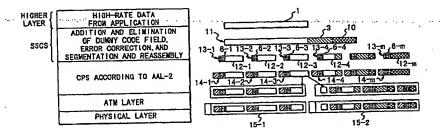


FIG. 1B



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### Description

### **TECHNICAL FIELD**

[0001] The present invention relates to a data communication method for transmitting high-rate data in addition to low-rate data, and to a data transmission processor and a data reception processor for realizing the data communication method. In the following, the low-rate data means low-volumes of data such as voice data, which may be transmitted at a low rate while the high-rate data means high volumes of data, which are preferably transmitted at a high-rate, other than voice data.

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### **TECHNICAL BACKGROUND**

[0002] Since traffics for not only voice data but also for other data have been rapidly increased in mobile communications, it is an important consideration to establish a large capacity mobile communication network. Moreover, for a mobile multimedia communication in future, it is preferable that communicators can equivalently treat voice data and other data, and therefore multiplex traffic communications capable of equivalently handling the voice data traffic and other data traffic is required. Such a technique is also required for point-to-point communications.

[0003] At present, ITU-T (International Telecommunication Union-Telecommunication Standardization Sector) is standardising the ATM (Asynchronous Transfer Mode) network system for the above-mentioned multiplex traffic communications. As a result, multimedia communications in the ATM networks including both of WANs and LANs can be realized. While the ATM network system is standardized, the standardization of the AAL (ATM Adaptation Layer) between the ATM layer and the higher layer is also carried out. Accordingly, the AAL type 2 protocol in accordance with service class B for transmitting coded data, such as coded voice data, at multiple data rate is standardized in order to lower transmission delay and to enhance transmission efficiency.

[0004] By the AAL type 2, low-rate data are transmitted through a plurality of channels provided with a single ATM virtual channel. Since mobile communications mainly treats voice traffics which are send at low and multiple data rate, the AAL type 2 is superior to be used in mobile communications. Therefore, it is preferable that the AAL type 2 is quickly introduced into mobile communications.

[0005] However, by the AAL type 2, it is impossible to transmit user data units having a length of 45 octets or more (one octet equals to eight bits). Therefore, if it is applied to multimedia communications in which high volumes of data, which should be usually sent at 64 kilobits per second or more, in addition to low volumes of data such as voice data, transmission rate of voice data

is delayed because of multiplexing with high volumes of data, so that talking quality may be decayed.

[0006] Instead, it would be possible to use the AAL type 5, which has been already standardized for high-rate data transmission. In this case, a statistical multiplex effect can be obtained because cells for voice data and high-rate data are integrally transmitted at the ATM layer and the lower layer. However, the statistical multiplex effect is limited at the ATM layer and the lower layer, and therefore, the transmission delay of voice data is not always prevented.

[0007] Therefore, to realize high-rate data transmission using with the AAL type 2, an improvement is proposed, in which the higher and extended sublayer of the AAL type 2 is equivalent to the AAL type 5 (Goran Eneroth at Ericsson, "Segmentation proposal for AAL-CU", ITU-T Study Group 13 Question 6 Madrid rapporteur meeting, Nov., 1996). Figure 8 shows usual hierarchical structure, including the AAL type 2, according to B-ISDN (Broadband aspects of ISDN). As shown in Figure 8, the high-rate data side of AAL type 2 is divided into two sublayers: the upper one is the service specific convergence sublayer (hereafter abbreviated as SSCS), and the lower one is the common part convergence sublayer (hereafter abbreviated as CPS).

By the above-mentioned proposal, the highrate data side of the layer structure shown in Figure 8 is modified into the structure in Figure 9A. That is, in the proposed structure, the SSCS is responsible for a function equivalent to the AAL type 5. Therefore, as shown in Figure 9A, when a user data unit 1 exceeding a fixed length (45 octets) is transferred to the SSCS of the AAL type 2 from the higher layer, a restorative code field 2 for error detection and a dummy code field 3 of variable length for padding (adjusting the length) are attached to the user data unit 1, whereby a data unit 4 is assembled in accordance with the AAL type 5 protocol. Next, the data unit 4 is divided into fixed-length segments 5-1 to 5-k (k is an integer equal to or more than 2). Successive information fields 6-1 to 6-k are then attached to the segments 5-1 to 5-k, whereby data blocks 7-1 to 7-k are prepared. Each of the successive code fields 6-1 to 6-k denotes that the corresponding segment is the last one or successive segment is following. The length of the dummy code field 3 is selected from a range between 0 and 44 octets so that the data length before dividing becomes integral multiples the fixed length, 45 octets. [0009] The data blocks 7-1 to 7-k are transferred to the CPS of the AAL type 2. In response, the data blocks 7-1 to 7-k are transformed into CPS packets 8-1 to 8-k by a process according to the AAL type 2 protocol, Next. the CPS packets 8-1 to 8-k are multiplexed and then transferred to the ATM layer. The multiplexed CPS packets are contained into ATM cells 9 in the ATM layer, and then the ATM cells 9 are transferred to a physical layer. [0010] By the above-described method, a high-rate data unit is first divided, and then multiplexed in the CPS. Therefore, it is expected that a statistical multiple

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effect can be also benefited at the CPS that is higher than the ATM layer.

However, by the function equivalent to AAL [0011] type 5 in the SSCS, unnecessary data elements, which are not used in the CPS, are included in the data blocks 7-1 to 7-k transferred to the CPS from the SSCS. As shown in Figure 9B, the restorative code field 2 includes an SSCS-UU (user-to-user) indication field of one octet, a CPI (common part indicator) field of one octet, a length indicator field of two octets, and an error correcting code field or cyclic redundancy check (CRC) field of four octets. However, the AAL type 2 does not necessarily utilize some of them. Consequently, the redundancy of transmitted data is high in comparison with the communications simply using with the AAL type 5 and the AAL type 2 in parallel, and transmission efficiency may be decreased

[0012] On the other hand, the layer higher than the AAL is utilized typically for managing voice communication and data communication. In voice communication, it is not always necessary that all of data elements be transmitted as long as a certain quality is secured, but the information should be sent in real time. On the other hand, in data communication, all of the data elements must be transmitted accurately although quickness is not needed very much. Therefore, if a data error is found, it is necessary to retransmit in data communication.

[0013] Thus, it is the function of the higher layer to decide the necessity of retransmission in despite that the lower layer is the AAL type 2 or not. When the retransmission is necessary, one or more data units (e.g., user frames) adapted to the higher layer are retransmitted from the higher layer. The user frame length adapted to the higher layer usually ranges between 256 octets and 64 kilo-octets.

[0014] On the other hand, the frame length should be small for transmission at the lower layers. For example, the frame length is 45 octets or less in accordance with the AAL type 2. Therefore, it is necessary to divide a user data frame from the higher layer into a plurality of segments before transferring them to the lower layers. However, since one or more user frames, which are longer, are retransmitted from the higher layer if a data error occurs, transmission efficiency may be decreased.

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### DISCLOSURE OF INVENTION

[0015] Accordingly, it is an object of the present invention to provide a data communication method for enhancing transmission efficiency, and to provide a data transmission processor and a data reception processor for realizing the method.

[0016] To solve the above problems, in one aspect of the present invention, the method comprises the steps of: dividing a data unit into a plurality of packets; attaching identifiers to the packets; transmitting the packets; receiving the packets; deciding whether all of the packets.

ets for reassembling the data unit are received in accordance with the identifiers; and reassembling the user data unit when the decision is affirmative.

### BRIEF DESCRIPTION OF DRAWINGS

### [0017]

Figure 1A is a conceptual diagram showing the hierarchical structure of the high-rate data side including the AAL type 2 according to B-ISDN and showing data formats respectively corresponding to the layers used for a data communication method according to a first embodiment of the present invention;

Figure 1B is an illustration showing the data structure of a restorative code field 10 in Figure 1A;

Figure 2 is an schematic illustration showing the structures of a transmission node 26 and a reception node 32 for realizing the data communication method using with the hierarchical structure in Figure 1A:

Figure 3 is a block diagram showing the structure of a high-rate data transmission processor 23 in the transmission node 26 in Figure 2 and the structure of a high-rate data reception processor 29 in the reception node 32 in Figure 2;

Figure 4A is a conceptual diagram showing the hierarchical structure of the high-rate data side including the AAL type 2 according to B-ISDN and showing data formats respectively corresponding to the layers used for a data communication method according to a second embodiment of the present invention;

Figure 4B is an illustration showing the data structure of restorative code field 40 in Figure 4B;

Figure 5 is an schematic illustration showing the structures of a transmission node 51 and a reception node 53 for realizing the data communication method using with the hierarchical structure in Figure 4A;

Figure 6 is a block diagram showing the structure of a high-rate data transmission processor 52 in the transmission node 51 in Figure 5 and the structure of a high-rate data reception processor 54 in the reception node 53 in Figure 5;

Figure 7 is a conceptual diagram for explaining a variation of the data communication method of the first embodiment;

Figure 8 is a diagram showing the hierarchical structure including the AAL type 2 according to B-ISDN;

Figures 9A is a conceptual diagram showing the hierarchical structure of the high-rate data side including the AAL type 2 according to B-ISDN and showing data formats respectively corresponding to the layers used for a communication system in which the higher and extended sublayer of the AAL

type 2 has a function equivalent to the AAL type 5; and

Figure 9B is an illustration showing the data structure of restorative code field 10 in Figure 9B.

### **BEST MODE FOR CARRYING OUT INVENTION**

[0018] Embodiments of the present invention will be described below with reference to the accompanying drawings. Although it is determined that the SSCS (service specific convergence sublayer) be served for high-rate data transmission in the recommendation about the AAL type 2 by ITU-T, functions of the SSCS are not determined specifically. Therefore, in the first and second embodiments of the present invention, the SSCS is responsible for new functions, so that low-rate data such as voice data and high-rate data can be transmitted with high quality.

### 1. FIRST EMBODIMENT

[0019] A first embodiment of the present invention will be described first with reference to Figures 1A to 3.

[0020] Figure 1A is a conceptual diagram showing the hierarchical structure of the high-rate data side including the AAL type 2 according to B-ISDN and showing data formats respectively corresponding to the layers used for a data communication method according to a first embodiment of the present invention. In Figure 1A, the same reference symbols are attached to elements that are common to Figure 9A. However, the functions in the SSCS in Figure 1A are limited to addition and elimination of dummy code fields, error correction, and segmentation and reassembly of data.

[0021] Moreover, the data format in Figure 1A is different from that in Figure 9A in that the restorative code field 10 is substituted for the restorative code field 2. Figure 1B is an illustration showing the data structure of the restorative code field 10. As shown in Figure 1B, the restorative code field 10 includes a length indicator field indicating the length of the user data unit 1, the length being possibly longer than 45 octets; and an error correcting code field for detecting deficiency and bit errors of the user data when restoring. Unlike the information field 2 of Figure 9B, the information field 10 does not include the SSCS-UU field and CPI field.

[0022] Since the restorative code field 2 is replaced by the restorative code field 10, new symbols are attached to some elements in Figure 1A: i.e., 11 to the data unit, 12-1 to 12-m to the segments, 13-1 to 13-m to the data blocks, 14-1 to 14-m to the CPS packets, 15 to the ATM cells, wherein m is an integer equal to or greater than 2, and equal to less than k (see Figure 9A). As similar to Figure 9A, the hierarchical structure in Figure 1A is a part of the entire hierarchical structure.

[0023] Figure 2 schematically shows the structures of a transmission node 26 and a reception node 32 for realizing data communication using with the above hier-

archical structure. In Figure 2, symbols 21-1 to 21-3 and 13-1 to 13-3 designate voice-data-user-terminals for realizing the higher layer of the voice data side according to B-ISDN. Each of the voice-data-user-terminals 21-1 to 21-3 and 13-1 to 13-3 comprises device drivers and stores applications or software in relation to voice data communication for preparing and relaying voice user data.

[0024] Symbols 22-1 to 22-3 and 30-1 to 30-3 denote high-rate-data-user-terminals for realizing the higher layer of the high-rate data side of the hierarchical structure according to B-ISDN. Each of the high-rate-data-user-terminal 22-1 to 22-3 and 30-1 to 30-3 comprises device drivers and stores applications or software in relation to high-rate communication for preparing and relaying user data. The number of each user-terminal may be optionally selected although Figure 2 shows three voice-data-user-terminals and three high-rate-data-user-terminals.

[0025] In the transmission node 26, each of high-rate data transmission processors 23-1 to 23-3 for realizing the SSCS, executes a process to user data units transferred from the corresponding high-rate-data-user-terminal 22-1, 22-2, or 22-3 as will be described later, and then transfers data blocks produced of the data units to a device for the layer lower than the SSCS. An AAL-2 converter 24 for realizing the CPS, attaches headers respectively to segments from any of the voice-datauser-terminals 21-1 to 21-3 and the high-rate data transmission processors 23, the header indicating data type, i.e., voice data or high-rate data. Then, the AAL-2 converter 24 converts the data blocks with the headers into CPS packets, multiplexes the CPS packets, and transfers them to a device for the lower layer. An AIF (ATM Interface) 25 for realizing the layer lower than the CPS, loads the CPS packets from the AAL-2 converter 24 in ATM cells, and then transmits the ATM cells through the ATM network.

[0026] In the reception node 32, an AIF 27 extracts the multiplexed CPS packets from the ATM cells transferred through the ATM network, and then conveys the CPS packets to the device for the higher layer. An AAL-2 converter 28 for realizing the CPS, divides the multiplexed CPS packets from the AIF 27 and transfers the data on the payloads of the CPS packets to a device for the higher layer (e.g., any of the voice-data-user-terminals 31-1 to 31-3) according to the type indicated by the header. Each of high-rate data reception processors 29-1 to 29-3 for realizing the SSCS, extracts data units from the received data blocks from the AAL-2 converter 28, and then transfers the user data units to the corresponding high-rate-data-user-terminal 30-1, 30-2, or 30-3 for the higher layer.

[0027] Next, structures of the high-rate data transmission processor 23 (any of 23-1 to 23-3) and the high-rate data reception processor 29 (any of 29-1 to 29-3) will be described below with reference to Figure 3.

[0028] In the high-rate data transmission processor

23, a dummy-information-field-attaching module 23-a prepares a dummy code field of a variable length (0 to 44 octets) and attaches it to the user data unit from the corresponding high-rate-data-user-terminal 22-1, 22-2, or 22-3 for the higher layer. The length of the dummy code field is adjusted so that the entire length of the data unit before dividing by a decomposing module 23-c is integral multiples the 45 octets.

[0029] An error-control-calculating module 23-b calculates an error correcting code (CRC) in accordance with a predetermined polynomial expression, referring to the user data unit with the dummy code field, and measures the length of the user data unit with the dummy code field. Then, the error control-calculating module 23-b attaches the error correcting code field and the length indicator field to the data unit. Consequently, the assembled data unit is constituted of the user data unit, dummy code field, length indicator field, and error-correcting code field, so that the error correcting code is served for only the user data unit, dummy code field, and length indicator field.

[0030] A decomposing module 23-c divides the data unit, constituted of the user data unit, dummy code field, length indicator field, and error correction code field, into fixed-length segments. Also, the decomposing module 23-c attaches successive code fields to the segments, respectively, and then transfers the segments to the AAL-2 converter 24 for the lower layer. The segments with the successive code fields are contained in payloads of the CPS packet. Each of the successive code fields denotes that the corresponding segment is the last one or successive segment is following as similar to the aforementioned successive code fields.

[0031] On the other hand, in the high-rate data reception processor 29, an assembling module 29-a temporarily stores the segments received from the AAL-2 converter 28, referring to the successive code fields of the segments, and reassembles a data unit from the segments.

An error control-calculating module 29-b [0032] receives the data unit assembled by the assembling module 29-a and detects the error correcting code of the data unit. In addition, the error control-calculating module 29-b calculates an error correcting code in accordance with the above-mentioned polynomial expression again. The existing error correcting code field is ignored in this calculation. Then, the error-control-calculating module 29-b compares the newly calculated code with the former code in the CRC field. When the comparison determination (error detection) is affirmative (the calculation code is equal to the former code), the error-control-calculating module 29-b transfers the data unit to a dummy-information-field-eliminating module 29-c. On the contrary, when the comparison determination is negative (the calculation result is not equal to the former CRC code), the error-control-calculating module 29-b deletes the data unit, and provides a notice with the dummy-information-field-eliminating module

29-c.

[0033] In response to reception of the data unit from the error-control-calculating module 29-b, the dummy-information-field-eliminating module 29-c extracts user data unit from the data unit in accordance with the length indicator of the data unit, and transfers the user data unit to the corresponding high-rate-data-user-terminal 30-1, 30-2, or 30-3 for the higher layer. On the contrary, if the notice of the negative comparison determination is received, the dummy-information-field-eliminating module 29-c carries out another operation that is predetermined for the higher layer, but the description of the operation is omitted.

[0034] Next, data transmission from the transmission node 26 to the reception node 32 will be described with reference to Figures 1A and 2. However, since voice data transmission from the voice-data-user-terminals 21-1 to 21-3 to the voice-data-user-terminals 31-1 to 31-3 is nothing more than the usual process for the AAL type 2, only high-rate data transmission process from the high-rate-data-user-terminals 22-1 to 22-3 to the high-rate-data-user-terminals 30-1 to 30-3 will be described below.

[0035] First, when the user data unit 1 having a length exceeding a fixed length is transferred from any of the high-rate-data-user-terminals 22-1 to 22-3 to the high-rate data transmission processor 23 in the transmission node 26, the dummy code field 3 is attached to the user data unit 1 in the high-rate data transmission processor 23 corresponding to the high-rate-data-user-terminal. Then, the restorative code field 10 is further attached to the user data unit 1, whereby the data unit 11 is constructed.

[0036] The data unit 11 is next divided into segments 12-1 to 12-m of a fixed-length (45 octets), and then, the successive code fields 6-1 to 6-m are attached to the segments, respectively, so that the data blocks 13-1 to 13-m are constructed. The data blocks 13-1 to 13-m are successively transferred to the AAL-2 converter 24, and are converted into CPS packets 14-1 to 14-m, and then multiplexed in the AAL-2 converter 24. The ATM cells 15 are then transmitted through the ATM network.

[0037] On the other hand, in the AIF 27 of the reception node 32, the ATM cells 15 supplied through the ATM network are received, and then the multiplexed CPS packets are extracted from the ATM cells 15. The multiplexed CPS packets are divided into the individual CPS packets 14-1 to 14-m in the AAL-2 converter 28. Moreover, in the AAL-2 converter 28, the data blocks 13-1 to 13-m are restored from the divided CPS packets 14-1 to 14-m, and successively transferred to the corresponding high-rate data reception processor 29. In the high-rate data reception processor 29, the data unit 11 is reassembled in accordance with the successive code fields 6-1 to 6-m in the data blocks 13-1 to 13-m.

[0038] In addition, the error detection of the combination of the user data unit 1, the dummy code field 3, and the length field is executed. That is, the recalculated

error correction code is compared with the former code in the data unit 11. When the comparison determination is effirmative, the user data unit 1 is extracted from the calculation in accordance with the length indicator of the restorative code field 10, and then the user data unit 1 is transferred to the corresponding high-rate-data-user-terminal. If the comparison determination is negative, the notice of the negative comparison determination is provided to the corresponding high-rate-data-user-terminal, and the data unit 11 is deleted.

[0039] As clearly understood from the above description, since the data unit 11 does not include the SSCS-UU field and CPI field, the data length of the unit 11 is two octets shorter than that of the data unit 4 (see Figures 9A and 9B). Therefore, this embodiment enables to efficiently transmit high-volumes of data while minimizing the transmission delay of voice data.

[0040] That is, by virtue of this embodiment, it is possible to transmit, using with the AAL type 2, not only low-rate data unit, such as voice data, of which the length is short, but also high-rate data unit, of which the length is greater than the payload length (45 octets) determined by the AAL type 2 protocol. Moreover, since the high-rate user data unit is divided into a plurality of segments and then the segments are transferred to the AAL type 2, it is possible to obtain the statistical multiple effect at the AAL type 2 and the lower layers. Therefore, it is possible to minimize transmission delay of low-rate data such as voice data. Moreover, because the length of the additional fields (restorative code field) for the user data unit 1 can be decreased, it is possible to prevent the transmission efficiency from being lowered.

### 2. SECOND EMBODIMENT

[0041] Next, a second embodiment of the present invention will be described below with reference to Figures 4A to 6.

[0042] Figure 4A is a conceptual diagram showing the hierarchical structure of the high-rate data side including the AAL type 2 according to B-ISDN and showing data formats respectively corresponding to the layers used for a data communication method according to a. second embodiment of the present invention. In Figure 4A, the same reference symbols are attached to elements that are common to Figure 1A. The hierarchical structure shown in Figure 4A is different from that in Figure 1A in that the SSCS in Figure 4A is not responsible for the addition and elimination of the dummy code field. Moreover, unlike the data format in Figure 1A, the dummy code field 3 is not attached and a restorative code field 40 is substituted for the restorative code field 10 in Figure 4A. Figure 4B is an illustration showing the data structure of the restorative code field 40. As shown in Figure 4B, the restorative code field 40 is constituted of only the error correcting code field for detecting deficiency and bit errors of the user data unit 1 when restoring and does not include the length indicator field

indicating the length of the user data unit 1.

[0044] Since the restorative code field 40 is substituted for the restorative code field 10 and the dummy code field 3 is excluded, new symbols are attached to some elements in Figure 4A: i.e., 41 to the data unit, 42-1 to 42-n to the segments, 43-1 to 43-n to the data blocks, 44-1 to 44-n to the CPS packets, 45 to the ATM cells, wherein n is an integer equal to or greater than 2, and equal to less than m (see Figure 1A). As similar to Figure 9A, the hierarchical structure in Figure 4A is a part of the entire hierarchical structure shown in Figure

[0045] Figure 5 is a schematic illustration showing the structures of a transmission node 51 and a reception node 53 for realizing the data communication method using with the hierarchical structure in Figure 4A. In Figure 5, the same reference symbols are attached to elements that are common to Figure 2.

[0046] In the transmission node 51 in Figure 5, each of high-rate data transmission processors 52-1 to 52-3 for realizing the SSCS, executes a process to user data units transferred from the corresponding high-rate-data-user-terminals 22-1, 22-2, or 22-3 as will be described later, and then transfers data blocks produced of the data units to a device for the layer lower than the SSCS. In the reception node 53, each of high-rate data reception processors 54-1 to 54-3 for realizing the SSCS, reassembles the user data unit from data blocks received from the AAL-2 converter 28, and then transfers the user data unit to one of the high-rate-data-user-terminals 30-1 to 30-3 for the higher layer.

[0047] Next, structures of the high-rate data transmission processor 52 (any of 52-1 to 52-3) and high-rate data reception processor 54 (any of 54-1 to 54-3) will be described below with reference to Figure 6.

[0048] In the high-rate data transmission processor 52, an error-control-calculating module 52-a receives the user data unit from the corresponding high-rate-data-user-terminals 22-1, 22-2, or 22-3. Then, the error-control-calculating module 52-a calculates an error correcting code (CRC) of four octets in accordance with a predetermined polynomial expression, referring to the user data unit, and then, attaches the error correcting code field to the user data unit. Consequently, the assembled data unit is constituted of the user data unit and error correcting code field, so that the error correcting code is served for only the user data unit.

[0049] A decomposing module 52-b divides the data unit, constituted of the user data unit and CRC field, into segments of a fixed length of 45 octets. However, it is possible that the last segment has a length less than the fixed length. Also, the decomposing module 52-b attaches successive code fields to the segments, respectively, and then transfers the segments to the AAL-2 converter 24 for the lower layer.

[0050] On the other hand, in the high-rate data reception processor 54, an assembling module 54-a temporarily stores the segments received from the AAL-2

converter 28, and reassembles a data unit from the segments.

An error control-calculating module 54-b receives the data unit assembled by the assembling module 54-a and detects the error correcting code of the data unit. In addition, the error control-calculating module 54-b calculates an error correcting code in accordance with the above-mentioned polynomial expression again. The existing CRC field attached to the data unit is ignored in this calculation, and therefore, the user data unit is only referred. Then, the error-control-calculating module 54-b compares the newly calculated code with the former code in the CRC field. When the comparison determination (error detection) is affirmative (the calculation code is equal to the former code), the error-control-calculating module 54-b transfers the data unit to the corresponding high-rate-data-user-terminal 30-1, 30-1, or 30-3. On the contrary, when the comparison determination is negative (the calculation result is not equal to the former CRC code), the errorcontrol-calculating module 54-b deletes the data unit, and provides a notice with the high-rate-data-user-terminal 30.

The data communication process from the 100521. transmission node to the reception node will be described below with reference to Figures 4A, 4B and 5. [0053]... First, when the user data unit 1 having a length exceeding the fixed length is transferred to the high-rate data transmission processor 52 from any of the highrate-data-user-terminals 22-1 to 22-3 in the transmission node 51, the restorative code field 40 is attached to the user data unit 1, whereby the data unit 41 is constructed in the high-rate data transmission processor 52 corresponding to the high-rate-data-user-terminal. The data unit 41 is then divided into the segments 42-1 to 42-p (wherein p = n-1) of the fixed length and segments 42-n of a length less than the fixed length. Next, the successive code fields 6-1 to 6-n are attached to the segments 42-1 to 42-n, respectively, whereby the data blocks 43-1 to 43-n are constructed. The data blocks 43-1 to 43-n are then transferred to the AAL-2 converter 24. Subsequent procedures in the transmission node 51 are the same as those of the first embodiment.

[0054] When the reception node 53 receives the ATM cells, the data blocks 43-1 to 43-n are restored by the AAL-2 converter 28 and are successively transferred to the corresponding high-rate data reception processor 54. Then, in the high-rate data reception processor 54, the data unit 41 is reassembled in accordance with the successive code fields 6-1 to 6-n in the data blocks 43-1 to 43-n. Next, the comparison determination (error detection) is carried out for the user data unit 1 in accordance with the error correcting code contained in the restorative code field 40 in the data unit 41. When the comparison determination is affirmative, the user data unit 1 is obtained by excluding the CRC field of four octets, and is transferred to the corresponding high-rate-data-user-terminal. However, if the comparison

determination is negative, a notice of the negative comparison determination is provided to the corresponding high-rate-data-user-terminal and the data unit 41 is deleted.

[0055] As described above, since the data unit 41 does not include the SSCS-UU field, CPI field, dummy code field, and length indicator field, the data length of the unit 41 is two through 46 octets shorter than that of the data unit 11. Therefore, this embodiment enables to further improve the advantages obtained by the first embodiment.

[0056] As an variation of the second embodiment, even if the restorative code field 40 includes the length indicator field, it is also possible to improve the advantages obtained by the first embodiment since the data unit 41 does not include the SSCS-UU field, CPI field, and dummy code field.

#### 3. THIRD EMBODIMENT

### 3.1. DATA FORMATS OF CPS PACKETS USED IN THIRD EMBODIMENT

[0057] A third embodiment will be described next. This embodiment is a modification of the first and second embodiment in order to conduct retransmission if a data error occurs.

[0058] First, data formats of CPS packets transmitted in the third embodiment will be described below. Figures 10 to 13 are conceptual diagrams showing data formats of CPS packets to which a retransmission control parameters are attached.

### 3.1.1. CPS PACKET FOR SENDING USUAL USER DATA

[0059] Figure 10 shows a CPS packet 610 for sending usual user data. The CPS packet 610 is constituted of a CPS packet header 612, a retransmission control code field 613, and a payload 625. In the retransmission control code field 613, contained are a information type identifier 616, a higher-layer frame sequence number 617, a CPS packet sequence number 618, and a parity bit 619 for the retransmission control code field 613.

[0060] The payload 625 is constituted of a user data field 614 and an error detecting code field 615 for the user data. The user data in the field 614 are data of a short higher-layer frame or a segment obtained by dividing a long higher-layer frame. The error detecting code in the field 615 is an error correcting code (CRC) for detecting an error of the user data field 614. However, when retransmission of the CPS packet 610, the error correcting code field 615 may be excluded from the retransmitted CPS packet 610.

[0061] The error detecting code attached by a source device is checked by a destination device, which receives the CPS packet 610. Accordingly, the destination device informs the source device of a request for

retransmission when the check result is extraordinary.

[0062] The information type identifier 616 of two bits identifies the type of the corresponding CPS packet:

the identifier 616 is "00" if the CPS packet is the CPS packet 610, for sending usual user data, which is sent by the source device to the destination device:

the identifier 616 is "01" if the CPS packet is a CPS packet 620 shown in Figure 11, for requesting confirmation, which is sent by the source device to the destination device;

the identifier 616 is "10" if the CPS packet is a CPS packet 630 shown in Figure 12, for confirmation, which is sent by the destination device to the source device; and

the identifier 616 is "11" if the CPS packet is a CPS packet 600 shown in Figure 13, for requesting retransmission, which is sent by the destination device to the source device.

[0063] Since Figure 10 shows the CPS packet 610, the identifier 616 is "00".

[0064] The higher-layer frame sequence number 617 indicates which higher-layer frame the user data in the payload 625 is originated from. More specifically, when transmitting every higher-layer frame, the sequence number 617 is changed, and then the changed number is attached to a CPS packet or all CPS packets originated from one higher-layer frame. Since the length of the number 617 is only two bits, the code is used in cyclic for the number 617. This number is given and managed by the source device.

[0065] The CPS packet sequence number 618 is used for supervising the order of succession of CPS packets in the destination device. When forming every CPS packet from the higher-layer frame, the sequence number 618 may be changed and then attached to the CPS packet. This is especially convenient if CPS packets are originated from segments obtained by dividing a higher-layer frame. This number is also given and managed by the source device. It is preferable that the code is used in cyclic for the number 618, too. It is also possible that, when transmitting a new higher-layer frame, the CPS packet sequence number 618 be reset.

### 3.1.2. CPS PACKET FOR REQUESTING CONFIRMA-TION

[0066] Figure 11 shows a CPS packet 620, for requesting confirmation, which may be transmitted from the source device to the destination device. In response to the CPS packet 620, the destination device may reply a confirmation to the source device. The CPS packet 620 includes the CPS packet header 612, which is the same as in the CPS packet 610. However, the information type identifier 616 is "01" in order to indicate that this packet 620 requests the confirmation. Although the

CPS packet 620 contains the higher-layer frame sequence number 617, CPS packet sequence number 618, and parity bit 619, the destination device does not use them. In addition, although a payload 626 is contained in the CPS packet 620 in Figure 11, it can be excluded from the CPS packet 620 since the destination device does not use it.

[0067] Further to the confirmation-requesting CPS packet 600, a timer may be equipped in the source device since it is possible that the destination device does not receive the confirmation-requesting CPS packet 620. In this case, the timer limit is set to a time period that is necessary for a confirmation CPS packet 630, which will be described later, to reach the source device. If the timer count is reached to the limit, the source device makes the retransmission.

#### 3.1.3. CONFIRMATION CPS PACKET

[0068] When the CPS packet 620 for requesting confirmation is transmitted from the source device to the destination device after the usual user data CPS packet 610 is normally received by the destination device, a CPS packet 630 for confirmation shown in Figure 12 is sent from the destination device to the source device. The CPS packet 630 includes the CPS packet header 612, which is the same as in the CPS packet 610. However, the information type identifier 616 is "10" in order to indicate that this packet 630 confirms the reception finish of sequential CPS packets 610 originated from a single higher-layer frame. Although the CPS packet 630 contains the higher-layer frame sequence number 617, CPS packet sequence number 618, and parity bit 619, the source device does not use them.

[0069] The CPS packet 630 contains a payload 633, wherein contained are a confirmation list field 631 and an error detection code field 632 for the confirmation list. The confirmation list in the field 631 enumerates CPS packets received by the destination device normally before sending the CPS packet 630 or in a certain time period before sending the CPS packet 630. The error detection code field 632 is a CRC for detecting an error of the list field 631, but is not necessarily contained in the CPS packet 630.

### 3.1.4. CPS PACKET FOR REQUESTING RETRANS-MISSION

[0070] When the CPS packet 620 for requesting confirmation is transmitted from the source device to the destination device after the usual user data CPS packet 610 is not appropriately received by the destination device, a CPS packet 600 for requesting retransmission shown in Figure 13 is transmitted from the destination device to the source device.

[0071] The CPS packet 600 includes the CPS packet header 612, which is the same as in the CPS packet 610. However, the information type identifier 616 is "11"

in order to indicate that this packet 600 requests the retransmission. Although the CPS packet 600 contains the higher-layer frame sequence number 617, CPS packet sequence number 618, and parity bit 619, the source device does not use them.

[0072] The CPS packet 600 contains a payload 603, wherein contained are a retransmission requesting list field 601 and error detection code field 602 for the retransmission requesting list. The retransmission requesting list in the field 631 enumerates CPS packets that are not received by the destination device normally before sending the CPS packet 600 or in a certain time period before sending the CPS packet 600. The error detection code field 602 is a CRC for detecting an error of the retransmission requesting list field 601, but is not necessarily contained in the CPS packet 600.

#### 3.2. SYSTEM STRUCTURE OF THIRD EMBODIMENT

#### 3.2.1. GENERIC STRUCTURE

[0073] Figure 14 shows the system structure according to the third embodiment. The system in Figure 14 comprises a source device 641 being equivalent to the transmission node 26 or 51 of the first or second embodiment; a higher-layer processor 642 being equivalent to one of the user-terminals 21-1 to 21-3 and 22-1 to 22-3 of the first or second embodiment; a destination device 648 being equivalent to the reception node 32 or 53 of the first or second embodiment; and a higher-layer processor 649 being equivalent to one of the user-terminals 30-1 to 30-3 and 31-1 to 31-3 of the first or second embodiment.

[0074] As shown in Figure 14, the devices 641 and 648 are of the same structure as each other in fact: i.e., modules 643 to 647 in the device 641 are the same as modules 6410 to 6414 in the device 648. Therefore, the same operation can be carried out by any of the devices 641 and 648. However, in order to simplify the description, the devices 641 and 648 are called source device and destination device, respectively.

[0075] The source device 641 comprises a dividing/reassembling module 643 for higher-layer frames. The dividing/reassembling module 643 receives a higher-layer frame supplied from the higher-layer processor 642 or 649 divides into segments if necessary. The dividing/reassembling module 643 then transfers CPS packet or packets to a retransmission-controlling module 644. On the other hand, when CPS packets are supplied from the retransmission-controlling module 644, the dividing/reassembling module 643 reassembles a higher-layer frame from the CPS packets if necessary. The module 643 then transfers the higher-layer frame to the higher-layer processor 642.

[0076] The retransmission-controlling module 644 transfers data between the dividing/reassembling module 643 and a CPS-packet-composing/decomposing module 645. In addition, the module 644 is responsible

for various functions for retransmission of CPS packets. The structure and functions of the retransmission-controlling module 644 will be described in detail later.

[0077] The CPS-packet-composing/decomposing module 645 attaches CPS packet header 612 to each of the non-divided higher-layer frame or divided segments, so as to compose CPS packets, and then transfers the CPS packets to an ATM-cell-composing/decomposing module 646. In reception mode, the CPS-packet-composing/decomposing module 645 eliminates the CPS packet header from each CPS packet received from the ATM-cell-composing/decomposing module 646, and then transfers the remaining data block to the retransmission-controlling module 644.

[0078] The ATM-cell-composing/decomposing module 646 attaches an ATM cell header to each of the CPS packets, so as to compose ATM cells, and then transfers the ATM cells to an ATM interface 647. In reception mode, the ATM-cell-composing/decomposing module 646 eliminates the ATM cell header from each ATM packet received from the ATM interface 647, and then transfers the remaining CPS packet to the CPS-packet-composing/decomposing module 645. The ATM interface 647 transmits and receives the ATM cells to and from the ATM interface 6414 of the destination device 648.

### 3.2.2. DIVIDING/REASSEMBLING MODULE 643 FOR HIGHER-LAYER

[0079] Figure 15 shows in detail the structure of the dividing/reassembling module 643 for higher-layer frames. The dividing/reassembling module 643 comprises a length-checking module 651 for higher-layer frames. The length-checking module 651 determines if higher-layer frames, transferred from the higher-layer processor 642, are longer than 45-X (octets) or not, wherein X is the length of the retransmission control code field 613. Specifically, in the embodiment, X is one octet (eight bits) as shown in Figures 10 to 13.

[0080] However, X may vary due to contents of the retransmission control code field 613. For example, it is possible that a CPS packet contains a division information field indicating if the CPS packet is originated from a segment divided from a higher-layer frame or not in order to facilitates reassembling the higher-layer frame. In this case, X is one octet longer than the length of the division information field.

[0081] If the length-checking module 651 determines that a higher-layer frame is longer than 45-X (octets), the checking module 651 transfers it to a multiple-packet-processing module 653. In the processing module 653, the higher-layer frame is divided into CPS packets. The CPS packets are then transferred to the retransmission-controlling module 644.

[0082] However, if a higher-layer frame is not longer than 45-X (octets), the frame is transferred through a single-packet-processing module 652 to the retransmis-

sion-controlling module 644.

[0083] The dividing/reassembling module 643 also includes a higher-layer-frame-sequence-number-managing module 654. When a new higher-layer is supplied to the length-checking module 651, the managing module 654 renews a higher-layer-frame-sequence-number signal 655.

### 3.2.3. TRANSMITTING ORGANIZATION OF RETRANSMISSION-CONTROLLING MODULE 644

[0084] Figures 16 and 17 show in detail the retransmission-controlling module 644. The retransmission-controlling module 644 comprises a transmitting organization shown in Figure 16. In the transmitting organization, a sequence-number-managing module 664 includes a higher-layer-frame-sequence-number-managing module 665 and a CPS-packet-sequence-number-managing module 666.

[0085] When a CPS packet is supplied to a retransmission-information-attaching module 662 in the transmitting organization of the retransmission-controlling module 644, the higher-layer-frame-sequence-number-managing module 665 inquires the higher-layer-frame-sequence-number-managing module 654 about the higher-layer frame sequence number corresponding to the supplied frame or segment. On the inquiry, the managing module 665 in the retransmission-controlling module 644 receives a higher-layer-frame-sequence-number signal 655 from the managing module 654 in the dividing/reassembling module 643, and then decides the higher-layer frame sequence number 617 in accordance with the received signal 655.

[0086] On the other hand, the CPS packet numbermanaging module 666 changes the CPS packet sequence number 618 when every frame or segment is received in the attaching module 662. [0087] The higher-layer frame sequence number 617

and the CPS packet sequence number 618 thus obtained are supplied to the retransmission-information-attaching module 662. The attaching module 662 attaches the sequence numbers 617 and 618 to each CPS packet received from the packet-processing module 652 or 653 in the dividing/reassembling module 643. [0088] The transmitting organization of the dividing/reassembling module 643 also includes a confirmation-requesting-CPS-packet-preparing module 661 that

tion-requesting-CPS-packet-preparing module 661 that prepares the CPS packets 620 for requesting confirmation at necessary occasions. That is, the CPS packets 620 are prepared periodically or in response to trigger signals independently of transmission of the higher-layer frame.

[0089] The transmitting organization of the dividing/reassembling module 643 also includes an error-detecting-code-attaching module 663 that attaches the error detecting code field 615 to each CPS packet. In addition, the module 663 attaches the retransmission control code field 613 to each of the CPS packets. In

case of the CPS packet 610 for sending usual user data produced from a higher-layer frame, the attaching module 663 gives "00" for the information type identifier 616 in the retransmission control code field 613 (see Figure 11).

[0090] In order to realize retransmission in response to the request, the CPS packets processed by the error-detecting-code-attaching module 663 are stored in a buffer 667. The buffer 667 for retransmission is connected with a retransmission-managing module 668. When the confirmation list data in the field 631 is supplied to the retransmission-managing module 668, the module 668 deletes the CPS packet or packets indicated in the list. On the other hand, when the retransmission requesting list data in the field 601 is supplied to the retransmission-managing module 668, the managing module 668 read out the CPS packet indicated in the list, and delivers the packet to the error-detecting-code-attaching module 663, whereby the packet, which could not be received normally, can be retransmitted.

### 3.2.4. RECEIVING ORGANIZATION OF RETRANS-MISSION-CONTROLLING MODULE 644

Furthermore, the retransmission-controlling [0091] module 644 comprises a receiving organization shown in Figure 17. As shown in Figure 17, the receiving organization includes a parity-bit-calculating module 676 for the retransmission control code field 613. The calculating module 676 calculates a number in the same manner for calculating the parity bit 619 with reference to the information type identifier 616, higherlayer frame sequence number 617, and CPS packet sequence number 618 in a CPS packet supplied from the CPS-packet-composing/decomposing module 645. Then, the calculating module 676 compares the parity bit 619 contained in the CPS packet with the newly calculated number. If they are not equal to each other, the calculating module 676 deletes the corresponding CPS packet.

[0092] On the other hand, if the calculated number is equal to the former parity bit 619, the corresponding CPS packet is transferred to an error-detecting-code-calculating module 675 for the user data field 614. The calculating module 675 calculates a code in the same manner for calculating the error detecting code field 615 with reference to the user data field 614. Then, the calculating module 675 compares the error detecting code contained in the field 615 of the CPS packet with the newly calculated code. If they are not equal to each other, the calculating module 675 provides an abnormal-CPS-packet-supervising module 672 with the retransmission control code in the field 613 of the corresponding CPS packet, and deletes this CPS packet.

[0093] On the other hand, when the calculated number is equal to the former error detecting code, the calculating module 675 provides a detection module 674 with the corresponding CPS packet. The detection

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module 691.

module 674 detects the information type identifier 616 in the CPS packet. When the detection module 674 for the information type identifier detects that the identifier 616 is "00", identifying the CPS packet 610 for sending user data, the packet is transferred to a retransmission-control-information-eliminating module 673. However, if the identifier 616 is "11" or "10", identifying the retransmission requesting CPS packet 600 or the confirmation CPS packet 630, the packet is transferred to the retransmission-controlling module 677.

[0094] The retransmission-controlling module 677 extracts the retransmission requesting list field 601 and the confirmation list field 631 and transfers the data of these lists to the retransmission-managing module 668 shown in Figure 16.

The usual user data CPS packet or packets, [0095] supplied to the retransmission-control-information-eliminating module 673, are stored in a buffer 678. The buffer 678 stores the CPS packet or packets until the reception is completed. In addition, the retransmission control code in the field 613 of the CPS packet is supplied to a normal-CPS-packet-supervising module 671, which supervises the higher-layer frame sequence number 617 and CPS packet sequence number 618. More specifically, the packet supervising module 671 determines that the order of succession of the supplied higher-layer frame sequence number 617 and CPS packet sequence number 618 is regular or not. If a lost sequence number 617 and/or 618 is detected, the normal-CPS-packet-supervising module 671 informs the abnormal-CPS-packet-supervising module 672 of the lost number.

[0096] When it is decided by the normal CPS-packet supervising module 671 that one or more CPS packets originated from one higher-layer frame are normally received, the retransmission-control-information-eliminating module 673 reads out the CPS packet or the CPS packets from the buffer 678, and supplies the packet or packets to the dividing/reassembling module 643 for higher-layer frames.

[0097] Otherwise, the abnormal-CPS-packet-supervising module 672 instructs the CPS-packet-composing/decomposing module 645 to send the CPS packet 600 for requesting retransmission of the CPS packet 610, in which the user data field 614 has an error, or which is deleted for a communication trouble. More specifically, the lost sequence numbers 617 and 618 of the CPS packet are provided with the composing/decomposing module 645.

[0098] On the contrary, if the normal-CPS-packet-supervising module 671 determines that the CPS packet or packets are normally received, the supervising module 671 instructs the CPS-packet-composing/decomposing module 645 to send the CPS packet 630 for confirmation. In this case, the received sequence numbers 617 and 618 of the CPS packet are provided with the composing/decomposing module 645.

3.2.5. CPS-PACKET-COMPOSING/DECOMPOSING MODULE 645

[0099] Figure 18 shows in detail the CPS-packet-composing/decomposing module 645. The module 645 comprises a CPS-packet-header-attaching/eliminating module 681 and a CPS-packet-header-checking module 682. The CPS-packet-header-attaching/eliminating module 681 attaches the CPS packet header 612 to a CPS packet supplied from the retransmission-controlling module 644, and then transfers it to the CPS-packet-header-checking module 682. In addition, the CPS-packet-header-attaching/eliminating module 681 eliminates the CPS packet header 612 from each CPS packet supplied from the CPS-packet-header-checking module 682, and then transfers the packet to the retransmission-controlling module 644.

[0100] The CPS-packet-header-checking module 682 determines if the CPS packet header is normal or not. If the header is normal, the checking module permits the corresponding packet to transfer, but otherwise, the checking module deletes the packet.

### 3.2.6. ATM-CELL-COMPOSING/DECOMPOSING MODULE 646

Figure 19 shows in detail the ATM-cell-com-[0101] posing/decomposing module 646. The module 646 comprises a CPS-packet-multiplexing/separating module 691, an ATM-cell-header-attaching/eliminating module 692, and an ATM-cell-header-checking module 693. The CPS-packet-multiplexing/separating module 691 multiplexes the CPS packets supplied from the module 645 into payloads of ATM cells, and then transfers the ATM cells to the ATM-cell-header-attaching/eliminating module 692. On the contrary, the module 691 extracts CPS packets from a payload of an ATM cell supplied from the module 692, and transfers the CPS packets to the CPS-packet-composing/decomposing module 645. The ATM-cell-header-attaching/eliminating [0102] module 692 attaches an ATM cell header to the payload of the ATM cell supplied from the CPS-packet-multiplexing/separating module 691, and then transfers the cell to the ATM-cell-header-checking module 693. In addition, the ATM-cell-header-attaching/eliminating module 692 eliminates the ATM cell header from an ATM cell supplied from the ATM-cell-header-checking module 693, and then transfers the remaining payload of the ATM cell to the CPS-packet-multiplexing/separating

[0103] The ATM-cell-header-checking module 693 determines if the ATM cell header is normal or not. If the header is normal, the checking module 693 permits the corresponding ATM cell to transfer, but otherwise, the checking module 693 deletes the ATM cell.

### 3.3. OPERATIONS OF THIRD EMBODIMENT

### 3.3.1. NORMAL COMMUNICATION OPERATION

[0104] When a higher-layer frame, prepared by the higher-layer processor 642, is supplied to the dividing/reassembling module 643, the higher-layer frame is divided into segments and then transferred to the retransmission-controlling module 644. Alternatively, the higher-layer frame is, as it is, transferred to the retransmission-controlling module 644 if it is short. The retransmission-controlling-module-644-attaches the retransmission control code field 613 to the frame or segment. The CPS-packet-composing/decomposing module 645 also attaches the CPS packet header 612 to the frame or segment, thereby preparing a CPS packet.

[0105] Next, the ATM-cell-composing/decomposing module 646 transforms the CPS packet or packets to an ATM cell, and then transfers the ATM cell to the ATM interface 647, so that the ATM cells are transmitted to the ATM interface 6414 of the destination device 648. Then, the CPS-packet-composing/decomposing module 6412 extracts the CPS packet or packets from the ATM cell, and the ATM-cell-composing/decomposing module 6413 deletes the CPS packet header 612 from each CPS packet:

[0106] If the destination device 648 receives normally, the retransmission-controlling module 6411 of the destination device 648 recognizes the frame or segments produced from one higher-layer frame. The dividing/reassembling module 6410 reassembles the higher-layer frame based on the frame or segments. The reassembled higher-layer frame is transferred to the higher-layer processor 649.

[0107] In addition, when the frame or segments produced from one higher-layer frame are recognized, the normal-CPS-packet-supervising module 671 in retransmission-controlling module 6411 instructs the CPSpacket-composing/decomposing module 645 to send the CPS packet 630 for confirmation. In response to the reception of the CPS packet or packets 630 by the source device 641, the reception notice is provided with the retransmission-managing module 668 of the source device 641, whereby the CPS packets stored in the buffer 667 for retransmission are deleted. Consequently, if all the CPS packets corresponding to a higher-layer frame are transmitted from the source device 641 to the destination device 648 normally, the contents of the buffer 667 for retransmission is erased. Therefore, the source device 641 recognizes that the transmission has ended normally.

### 3.3.2. OPERATION WHEN COMMUNICATION TROU-BLE (USER DATA ERROR) OCCURS

[0108] However, if the error-detecting-code-calculating module 675 in the retransmission-controlling mod-

ule 6411 (Figure 17) detects a contradiction of the user data field 614 and the error detecting code in the field 615, the higher-layer frame sequence number 617 and the CPS packet sequence number 618 of the CPS packet is provided with the abnormal-CPS-packet-supervising module 672, and the CPS packet is deleted by the retransmission-controlling module 6411.

[0109] The abnormal-CPS-packet-supervising module 672 then instructs the CPS-packet-composing/decomposing module 645 to send the retransmission requesting CPS packet 600 for the CPS-packet, of which the error has be detected. In response, the destination device 648 transmits the CPS packet 600 to the source device 641 whereby the source device 641 makes the retransmission.

### 3.3.3. OPERATION WHEN COMMUNICATION TROUBLE (CPS PACKET LOSING) OCCURS

[0110] When an ATM cell header is determined to be abnormal by the ATM-cell-header-checking module 693, the corresponding ATM cell is also deleted. Furthermore, when a CPS packet header 612 is determined to be abnormal by the CPS-packet-header-checking module 682 or a parity bit 619 is determined to be abnormal by the parity-bit-calculating module 676, the corresponding CPS packet is also deleted. Consequently, in these cases, the CPS packet does not reach the error-detecting-code-calculating module 675, and the abnormal-CPS-packet-supervising module 672 cannot recognize the communication error quickly.

[0111] However, if a CPS packet is lost, the normal-CPS-packet-supervising module 671 determines that order of succession of the supplied higher-layer frame sequence number 617 and CPS packet sequence number 618 is irregular. Therefore, the normal-CPS-packet-supervising module 671 informs the abnormal-CPS-packet-supervising module 672 of the lost number, so that the CPS packet, to which a communication trouble occurred, can be specified. After that, the operation is similar to the above-mentioned error detection of the user data field 614. That is, the destination device 648 transmits the retransmission requesting CPS packet 600 to the source device 641, whereby the source device 641 makes the retransmission.

### 4. VARIATIONS AND MODIFICATIONS

[0112] Various variations and modifications of the present invention can be made without departing from the spirit or essential characteristics thereof. Therefore, the above embodiments are only examples and thus, the scope of invention is not restrictively defined by the above embodiments. The scope of the present invention is determined by the attached claims and rather than by the description. Moreover, all variations and modifications equivalent to the characteristics in the claims are included in the scope of the present inventions.

tion.

[0113] For example, the above embodiments can be modified as described below.

[0114] The length of the restorative code field, the subfields therein, and the segments can be optional as long as the embodiments accomplish the advantages.

[0115] The functional structure shown in Figure 2 or 5 can be realized by any of software and hardware.

[0116] It is possible to realize the functions of a highrate data transmission (reception) processor, AAL type-2 converter, and AIF by different nodes in a LAN.

[0117] It is also possible to make a single unit transmission/reception node by combining the transmission node 26 with the reception node 32 in Figure 2 or the transmission node 51 with the reception node 53 in Figure 5.

[0118] In the embodiments, the SSCS in Figure 1A executes the addition and elimination of dummy code fields, error correction, and segmentation and reassembly of data, and the SSCS in Figure 4A executes the error correction and segmentation and reassembly of data. However, these functions can be carried out by one of layers higher than the SSCS.

[0119] Furthermore, it is possible that the successive code attached by the SSCS be not contained in the payload of the CPS packet, rather be contained in the CPS packet header. Figure 7 shows the data formats according to this variation. In Figure 7, the same reference symbols are attached to elements that are common to Figure 1A. In this variation, the location of the successive information is not limited as long as it is loaded in the CPS header, however, it is preferable to load the successive information into the UUI (User-to-User Indication) field of the CPS packet header.

[0120] Moreover, Figure 20 shows a further modification of the variation shown in Figure 7. In Figure 20, the same reference symbols are attached to elements that are common to Figure 7. In this modification, the restorative code field 10 is excluded, and user data unit 1 is transferred between the higher layer and the CPS of AAL type 2 directly. In the modification, it is possible to readily transmit data since the restorative code field 10 used for retransmission is excluded. This modification may be used for a higher-layer application, which does not make retransmission, according to, for example, unrestricted digital information transfer service. Figure 21 exemplifies the code format of the UUI of the CPS packet header.

[0121] In the third embodiment, the ATM-cell-composing/decomposing module 646 loads the CPS packets into the payloads of the ATM cells. At this time, if a payload of an ATM cell contains apart of another CPS packet, only a part of the next CPS packet can be contained into the ATM cell. In this case, if the ATM cell is lost, the source device 641 should retransmit a plurality of CPS packets corresponding to the ATM cell, so that the overhead time involved for retransmission will be increased. However, when the ATM cell is lost, it is pref-

erable to retransmit a small number of CPS packets. Accordingly, the third embodiment can be modified as described below.

### FIRST MODIFICATION OF THIRD EMBODIMENT

[0122] After a series of the CPS packets divided from one higher-layer frame are loaded into ATM cells, the ATM-cell-composing/decomposing module 646 informs the dividing/reassembling module 643 of a remaining space length of the payload of the last ATM cell. The dividing/reassembling module 643 determines a payload length of the top CPS packet, which is the beginning part of a next transferred higher-layer frame from the higher-layer processor 642, in order to load it into the remaining payload space of the last ATM cell. For example, the payload length Lc1 of the top CPS packet is determined to satisfy the next equation.

$$. Lc1 = La - (Loa + Lor)$$

where La is the length of the remaining payload space of the ATM cell, Loa is the length of the overhead for AAL type 2, and Lor is the length of the overhead for retransmission.

[0123] The dividing/reassembling module 643 then determines the payload length Lc2 (expressed by octets) of the succeeding CPS packets, which should be prepared from the same higher-layer frame, to a predetermined length to satisfy the next equation.

$$Lc2 = 48 - (Loa + Lor)$$

where 48 (octets) is the entire length of the payloads of ATM cells.

[0124] The divided CPS packets are successively loaded into the payloads of the ATM cells. Except for the top and end CPS packets prepared from one higher-layer frame, it is not allowed to multiplex each of the CPS packets with another CPS packet from another connected terminal in one ATM cell payload.

[0125] In this modification, each of the CPS packets (except for the top and end CPS packets prepared from one higher-layer frame) occupies one ATM cell payload. Therefore, although one ATM cell is lost, it is enough to retransmit only a single CPS packet corresponding to the lost ATM cell.

### SECOND MODIFICATION OF THIRD EMBODIMENT

[0126] The ATM-cell-composing/decomposing module 646 always informs the dividing/reassembling module 643 of the remaining payload length of the last ATM cell. On the basis of the information, the dividing/reassembling module 643 always determines a payload length of the next CPS packet in order to load it into the remaining payload space of the last ATM cell. Namely, the payload length Lc of the top CPS packets is deter-

mined to satisfy the next equation.

where La is the length of the remaining payload space of the ATM cell, Loa is the length of the overhead for AAL type 2, and Lor is the length of the overhead for retransmission.

[0127] In this modification, it is allowed to multiplex each of the CPS packets with another CPS packet only from another connected terminal in one ATM cell payload. Since succesive CPS packets divided from a single higher-layer frame are not contained in one ATM cell, it is enough to retransmit only a single CPS packet corresponding to a lost ATM cell.

### THIRD MODIFICATION OF THIRD EMBODIMENT

[0128] The ATM-cell-composing/decomposing module 646 always pads payloads of the ATM cells (puts unnecessary data into the payloads of the ATM cells). The dividing/reassembling module 643 always divides the long higher-layer frames into CPS packets in such a manner that the payload length Lc (expressed by octets) of the succeeding CPS packets is a predetermined length to satisfy the next equation.

$$Lc = 48 - (Loa + Lor)$$

where 48 (octets) is the entire length of the payloads of ATM cells, Loa is the length of the overhead for AAL type 2, and Lor is the length of the overhead for retransmission. The divided CPS packets are successively loaded into the payloads of the ATM cells, respectively while the ATM cell payloads are padded. It is also not allowed to multiplex each of the CPS packets with another CPS packet from another connected terminal in one ATM cell payload.

[0129] In this modification, each of the CPS packets occupies one ATM cell payload. Therefore, although one ATM cell is lost, it is enough to retransmit only a single CPS packet corresponding to the lost ATM cell.

### **Claims**

1. A data communication method for transmitting high-rate user data, having a length exceeding a predetermined length, using with AAL type 2, the AAL type 2 originally capable of transmitting user data units to an ATM network, each of the user data units having a length shorter than or equal to the predetermined length and being transferred from a layer that is higher than the AAL type 2, the AAL type 2 originally capable of extracting data from packets received from the ATM network and transferring the data to the higher layer, a sublayer being situated between the higher layer and the AAL type 2, the method comprising the steps of:

attaching a restorative code to a user data unit using with the sublayer, thereby preparing a data unit, the restorative code being served for restoring the user data unit, the restorative code being an error detection code or error correcting code for checking the user data unit;

dividing the data unit into segments having a length equal to or less than the predetermined length;

attaching successive codes to the segments, respectively, each of the successive codes indicating whether successive segment is following to the corresponding segment or not;

transferring the segments with the successive codes to the AAL type 2 in order to load the segments with the successive codes into packets:

reassembling the data unit from the segments received from the AAL type 2 in accordance with the successive codes attached to the segments;

and

restoring the user data unit in accordance with the restorative code in the data unit and transferring the user data unit to the higher layer.

- 2. A data communication method according to claim 1, wherein the sublayer divides the data unit into the segments, attaches the successive codes to the segments, respectively, and transfers the segments with the successive codes to the AAL type 2 in order to load the segments with the successive codes into packets when the sublayer is served for transmission; and the sublayer reassembles the data unit from the segments in accordance with the successive codes extracted from the packets when the sublayer is served for reception.
- A communication method according to claim 2, wherein the restorative code comprises a length indicator indicating a length of the user data unit and an error detection code or error correction code for checking the user data unit and the length indicator.
- 4. A communication method according to claim 2, wherein the sublayer attaches a dummy code and the restorative code to the user data unit, thereby preparing the data unit when the sublayer is served for transmission, a length of the dummy code being determined according to the user data unit, the restorative code being served for restoring the user data unit and the dummy code; and the sublayer restores the user data unit in accordance with the restorative code in the reassembled data unit and transfers the user data unit to the higher layer, the restorative code being constituted of a length indicator indicating a length of the user data unit and an

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error detection code or error correction code for checking the user data unit, dummy code and length indicator.

- A communication method according to claim 4, 5 wherein a length of the dummy code is integral multiples the predetermined length.
- 6. A communication method according to one of claims 1 to 4, wherein a last segment among the segments prepared from the data unit has a length less than the predetermined length.
- 7. A data transmission processor for transferring a data unit to AAL type 2, the AAL type 2 originally capable of transmitting user data units to an ATM network, each of the user data units having a length shorter than or equal to the predetermined length and being transferred from a layer that is higher than the AAL, type 2, the data transmission processor comprising:

transmission processing means for attaching a restorative code to a user data unit using with the sublayer, thereby preparing a data unit, the restorative code being served for restoring the user data unit, the restorative code being an error detection code or error correcting code for checking the user data unit; the transmission processing means dividing the data unit into segments having a length equal to or less than predetermined length; transmission processing means attaching successive codes to the segments, respectively, each of the successive codes indicating whether successive segment is following to the corresponding segment or not; and transmission processing means transferring the segments with the successive codes to the AAL type 2 in order to load the segments with the successive codes into packets.

- 8. A data transmission processor according to claim 7, wherein the transmission processing means divides the data unit into the segments, attaches the successive codes to the segments, respectively, and transfers the segments with the successive codes to the AAL type 2 in order to load the segments with the successive codes into packets.
- A data transmission processor according to claim 8, wherein the restorative code comprises a length indicator indicating a length of the user data unit and an error detection code or error correction code for checking the user data unit and the length indicator.
- 10. A data transmission processor according to claim 8,

wherein the transmission processing means attaches a dummy code and the restorative code to the user data unit, thereby preparing the data unit, a length of the dummy code being determined according to the user data unit, the restorative code being served for restoring the user data unit and the dummy code, the restorative code being constituted of a length indicator indicating a length of the user data unit and an error detection code or error correction code for checking the user data unit, dummy code and length indicator.

- 11. A data transmission processor according to claim 10, wherein the length of the dummy code is adjusted so that the length of the data unit is integral multiples the predetermined length.
- 12. A data transmission processor according to one of claims 7 to 10, wherein a last segment among the segments prepared from the data unit has a length less than the predetermined length.
- 13. A data reception processor using with AAL type 2, the AAL type 2 extracting data from packets received from an ATM network and transferring the data to a layer higher than the AAL type 2, the data reception processor restoring a high-rate user data unit from the data extracted by the AAL type 2 and transferring the user data unit to the higher layer, the data reception processor comprising:

reception processing means for reassembling a data unit from segments extracted from the packets in accordance with successive codes attached to the segments, and for restoring a user data unit in accordance with a restorative code in the data unit, the restorative code being an error detection code or error correcting code for checking the user data unit, each of the successive codes indicating whether successive segment is following to the corresponding segment or not.

- 14. A data reception processor according to claim 13, wherein the reception processing means reassembles a data unit from segments extracted from the packets in accordance with successive codes extracted from headers of the packets.
- 50 15. A data reception processor according to claim 14, wherein the restorative code is constituted of a length indicator indicating a length of the user data unit and an error detection code or error correction code for checking the user data unit and length indicator.
  - 16. A data reception processor according to claim 14, wherein the data unit contains a dummy code

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attached by a source device if necessary, the restorative code being constituted of a length indicator indicating a length of the user data unit and an error detection code or error correction code for checking the user data unit, dummy code and length indicator, the reception processing means reassembling the data unit in accordance with the length indicator.

- 17. A data reception processor according to claim 16, wherein a last segment among the segments prepared from the data unit has a length less than the predetermined length.
- 18. A data reception processor according to one of claims 13 to 17, wherein the reception processing means, using with the error correction code, conducts a detection of errors in the user data unit, and transfers an information based on the detection to the higher layer.
- 19. A data reception processor according to one of claims 13 to 17, wherein the reception processing means, using with the error correction code, conducts a detection of errors in the user data unit, and transfers an information based on the detection, without the user data unit, to the higher layer if the reception producing means detects a error in the user data unit.
- 20. A data reception processor according to one of claims 13 to 17, wherein the reception processing means, using with the error correction code, conducts a detection of errors in the user data unit, and transfers the user data unit to the higher layer if the reception producing means detects no errors in the user data unit.
- 21. A data communication method using with a lower layer, the lower layer originally capable of transmitting user data units to a network, each of the user data units having a length shorter than or equal to a predetermined length and being transferred from a layer that is higher than the lower layer, the lower layer originally capable of extracting data from packets received from the network and transferring the data to the higher layer, the method comprising the steps of:

providing a sublayer between the higher layer and the lower layer;

attaching a restorative code to a user data unit using with the sublayer, thereby preparing a data unit, the restorative code being served for restoring the user data unit, the restorative code being an error detection code or error correcting code for checking the user data unit; dividing the data unit into segments having a

length equal to or less than the predetermined length;

attaching successive codes to the segments, respectively, each of the successive codes indicating whether successive segment is following to the corresponding segment or not;

transferring the segments with the successive codes to the lower layer in order to load the segments with the successive codes into packets:

transferring the segments with the successive codes from the lower layer to the sublayer; reassembling the data unit from the segments received from the lower layer in accordance with the successive codes attached to the segments; and

restoring the user data unit in accordance with the restorative code in the data unit and transferring the user data unit to the higher layer.

22. A data communication method comprising the steps of:

dividing a data unit into a plurality of packets; attaching identifiers to the packets; transmitting the packets; receiving the packets; deciding whether all of the packets for reassembling the data unit are received in accordance with the identifiers; and reassembling the user data unit when the deci-

23. A data communication method according to claim 22, wherein the identifiers are numbers which are changed in cyclic.

sion is affirmative.

- 24. A data communication method according to claim 23, wherein it is decided whether the identifiers of the received packets changes in cyclic or not in the decision step, the method further comprising a step of outputting a retransmission requesting signal indicating the identifier if the decision is negative.
- 25. A data communication method according to claim 24, wherein each of the identifiers comprise a first value which changes due to every packet and a second value which changes due to every user data unit
- 26. A data transmission method comprising the steps of:

dividing a data unit into a plurality of packets; attaching identifiers to the packets; transmitting the packets; receiving the packets; conducting a first decision for deciding if the received packets are normal or not in accordance with contents of the packets;

conducting a second decision for deciding if there is a lost packet or not in accordance with the identifiers attached to the packets if the first 5 decision is affirmative; and

outputting a retransmission requesting signal indicating an identifier of an abnormal packet or a lost packet if the first decision is negative or if it is decided that there is the lost packet in the 10 second decision.

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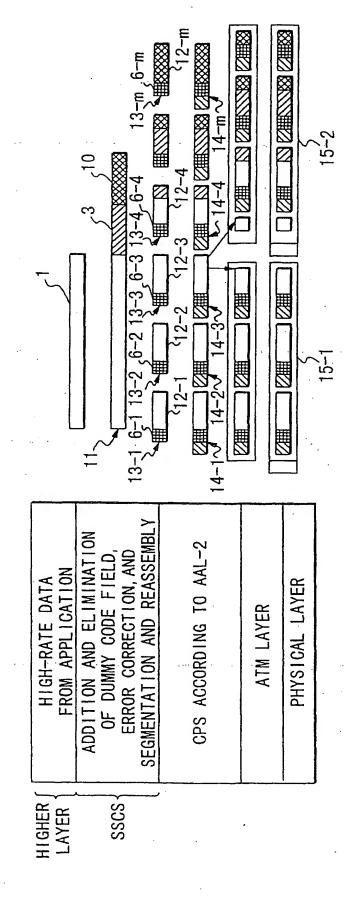
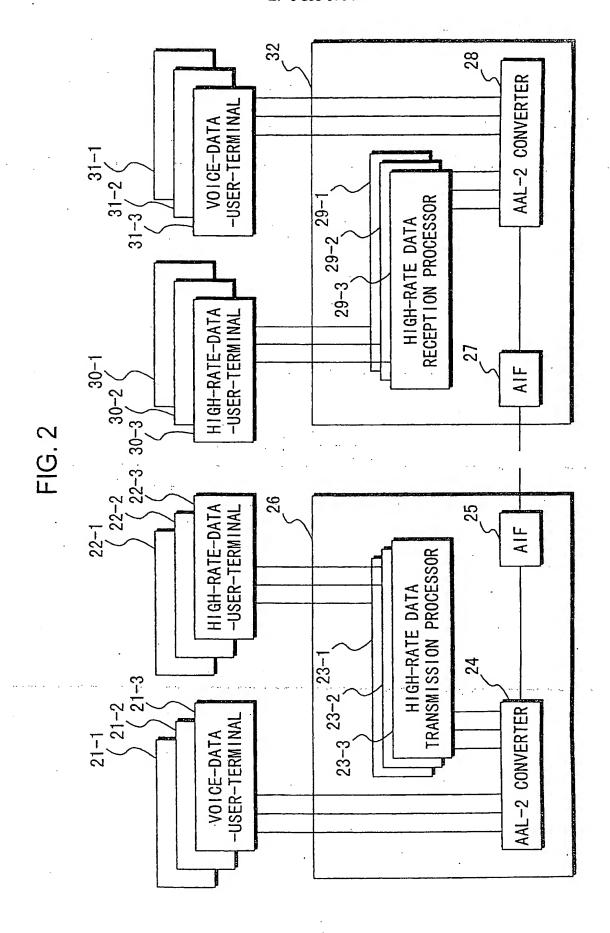
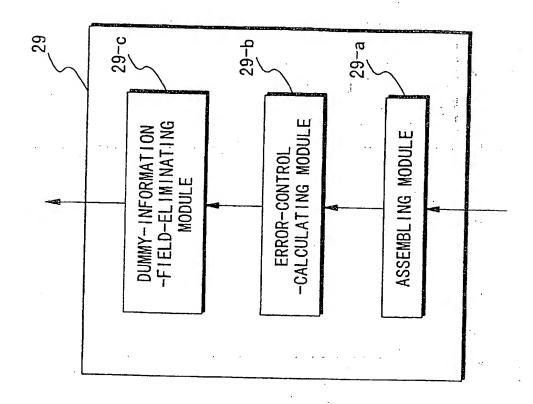
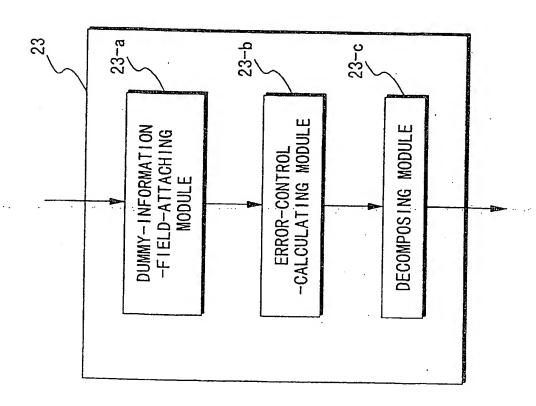


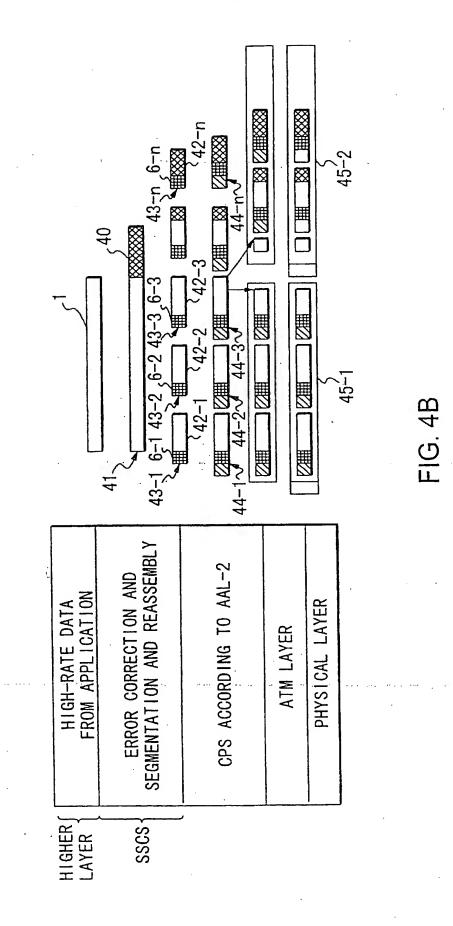
FIG. 1B



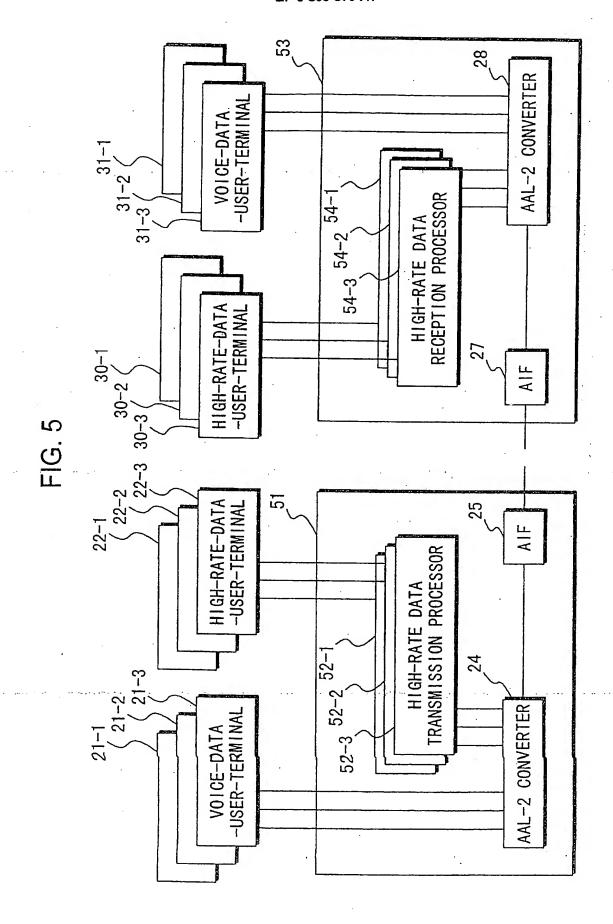


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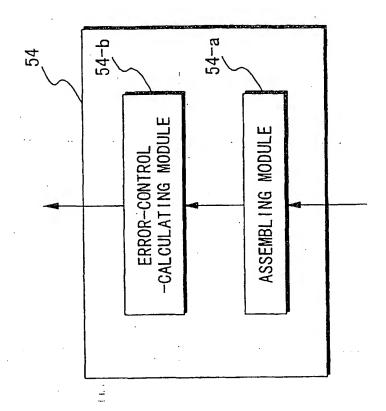
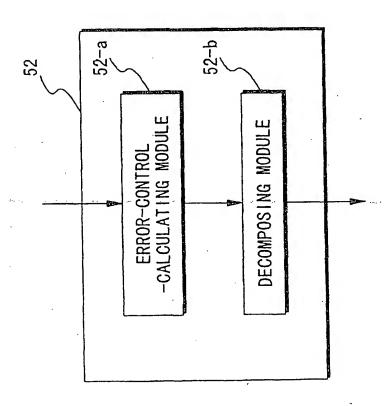
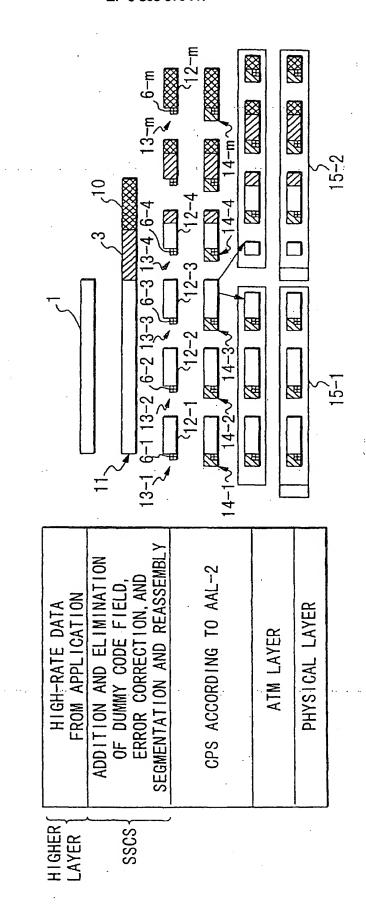
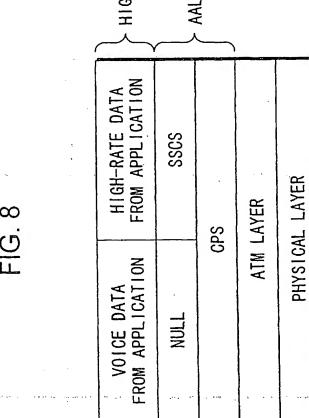
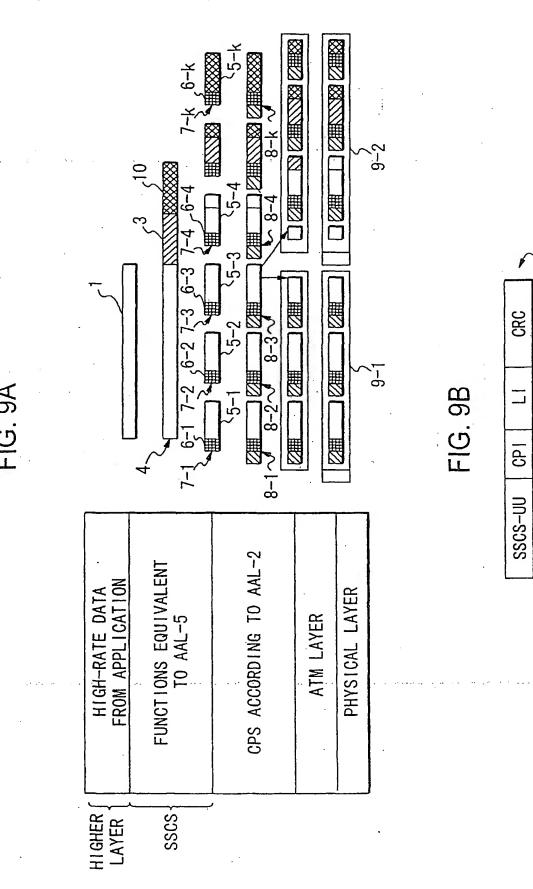


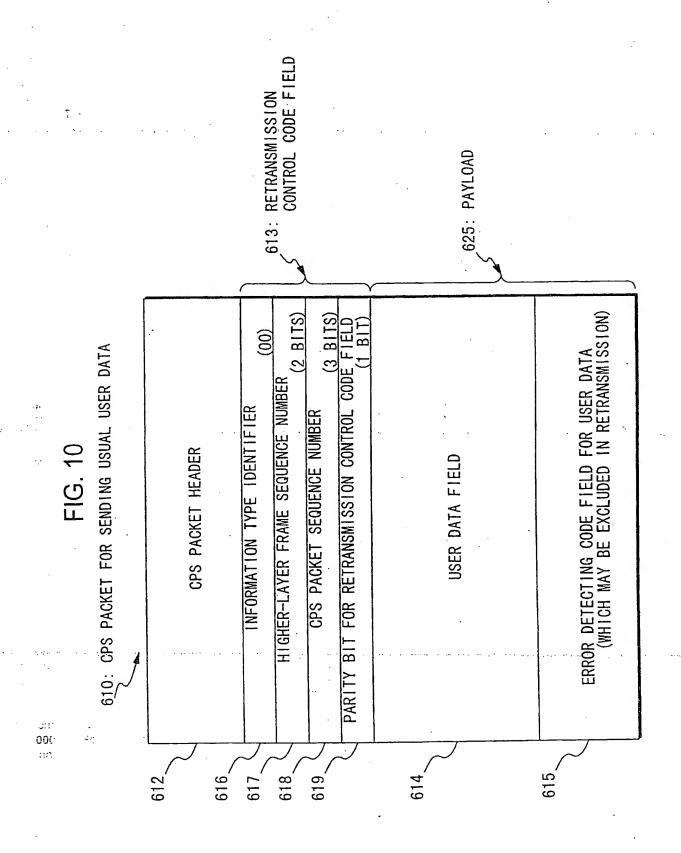
FIG. 6

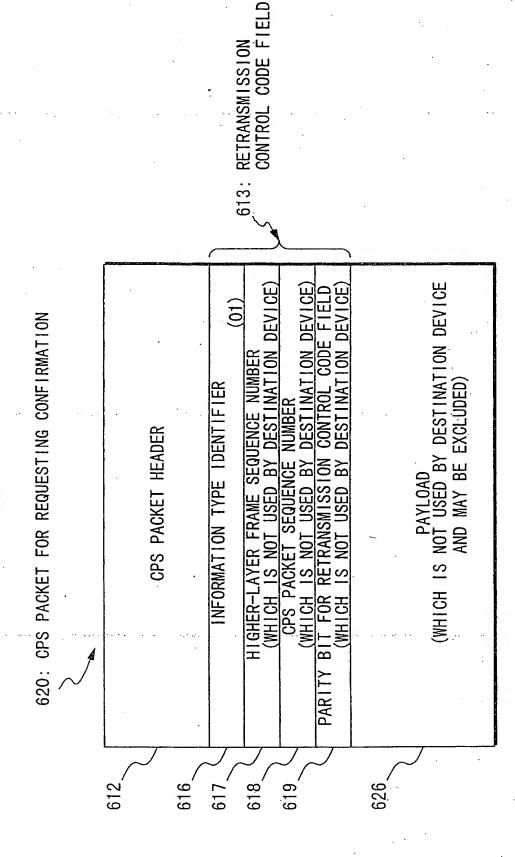


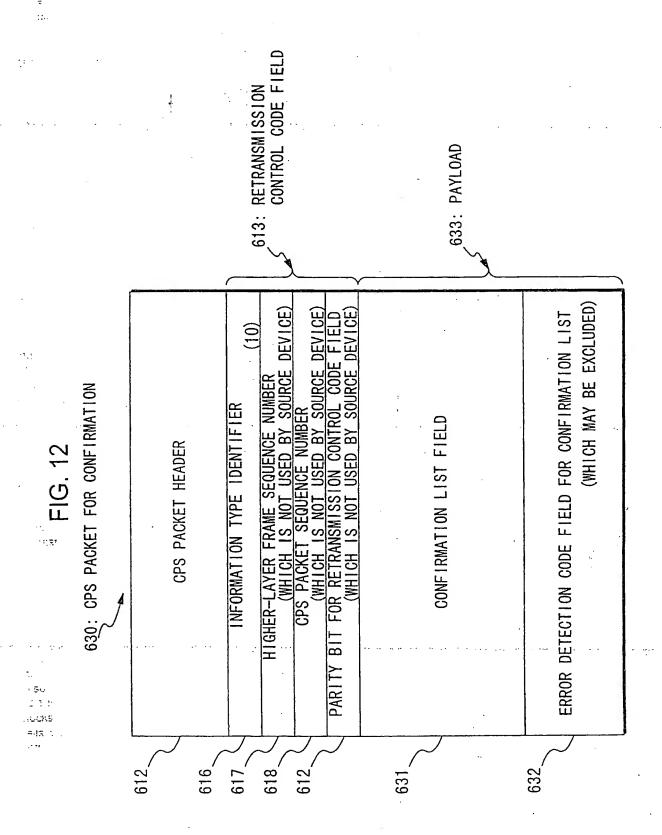


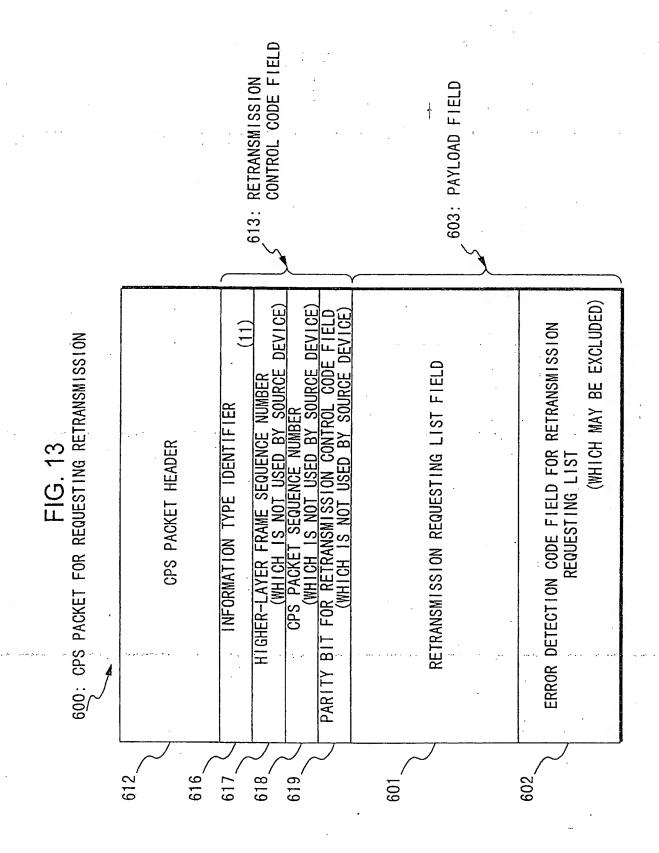


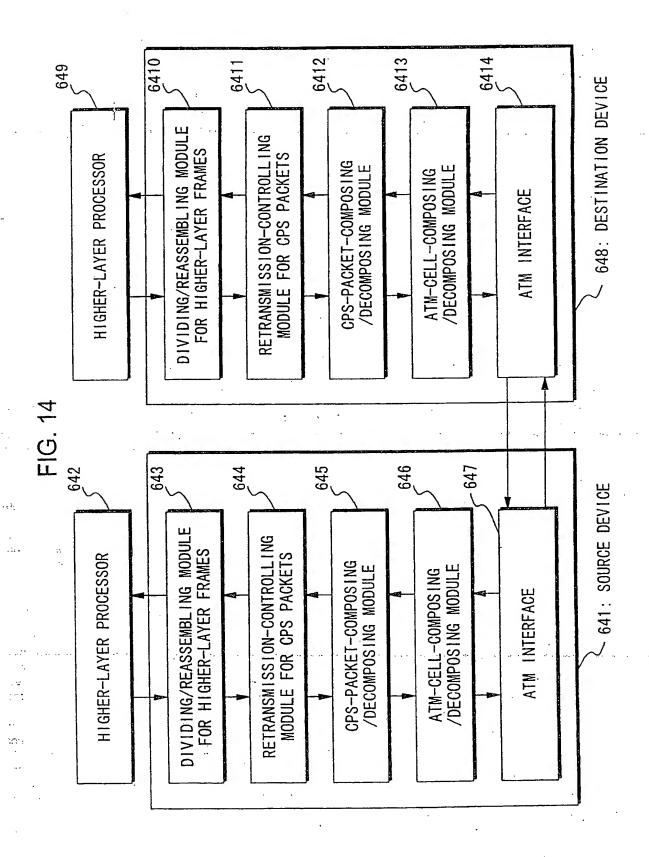












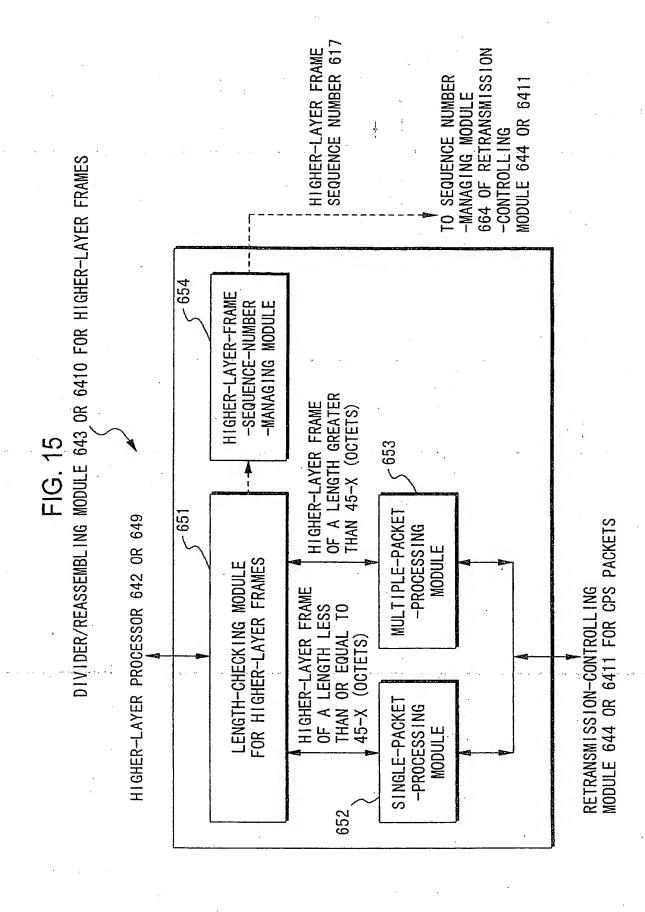
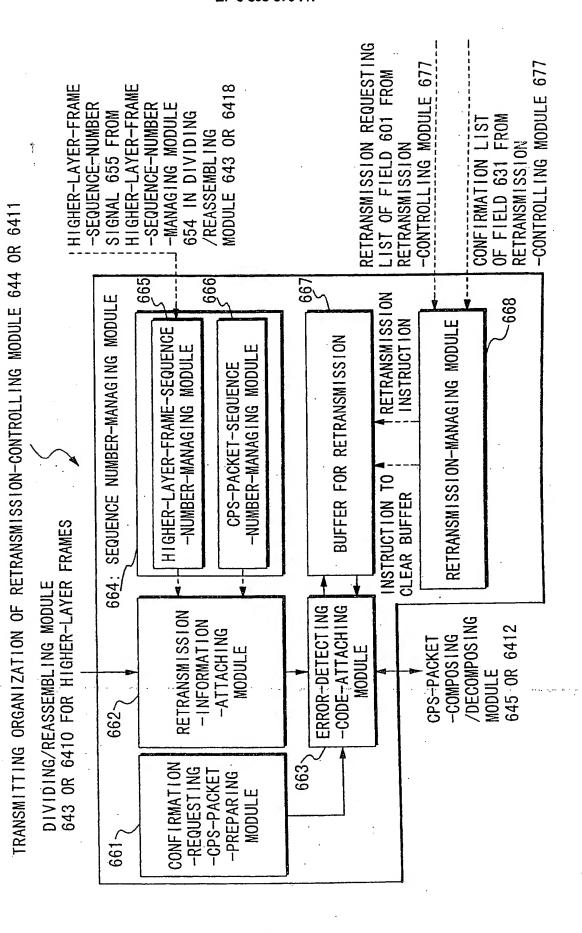
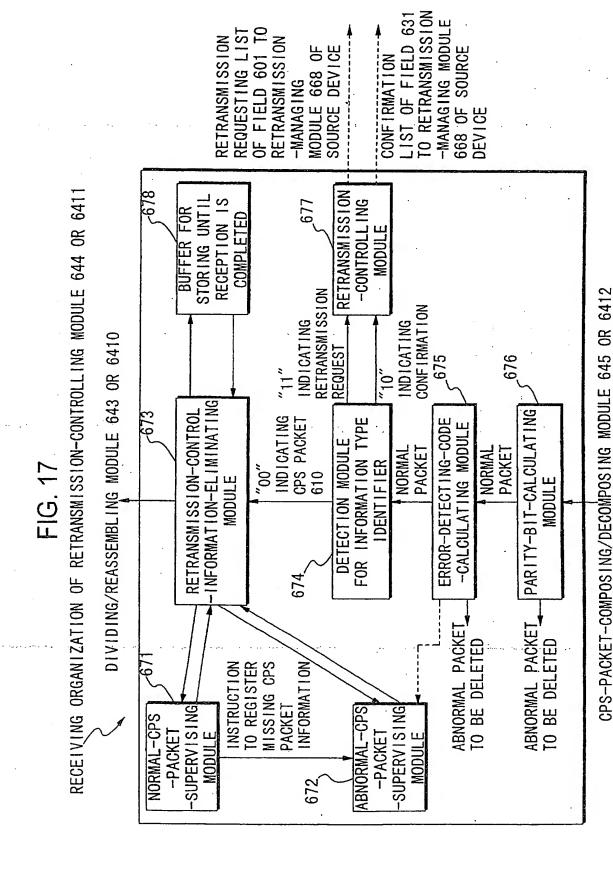
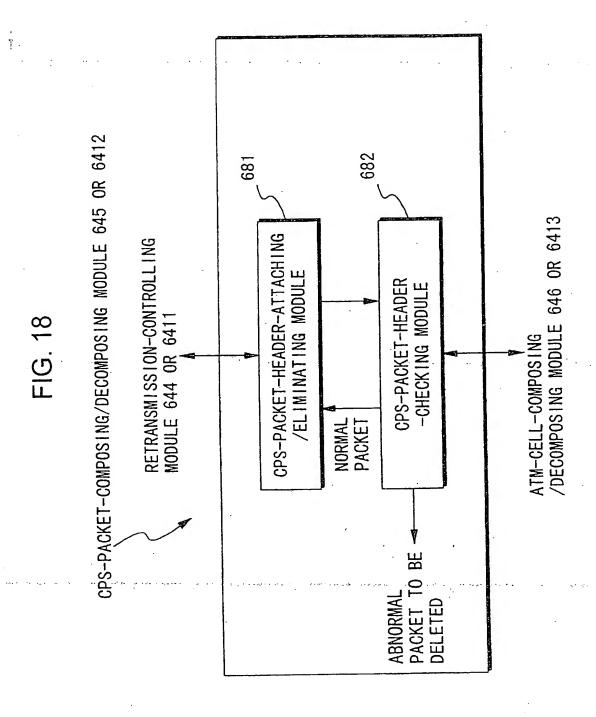
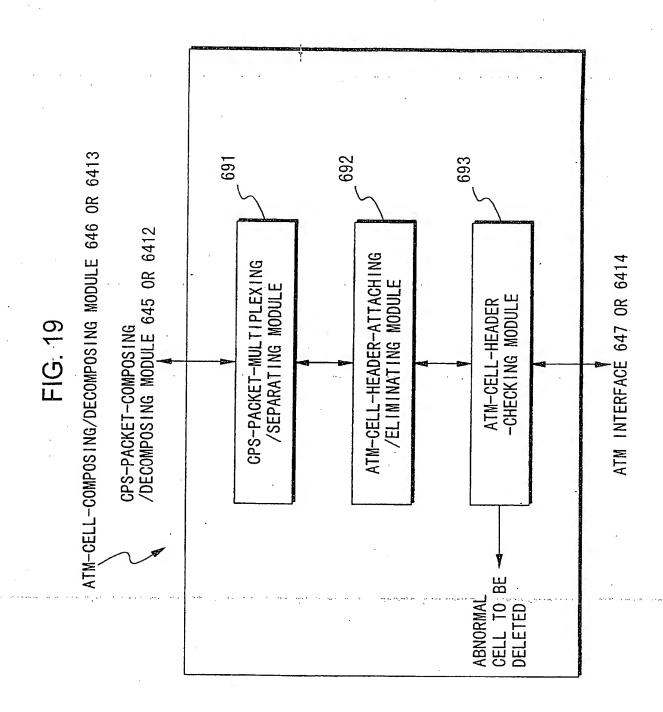


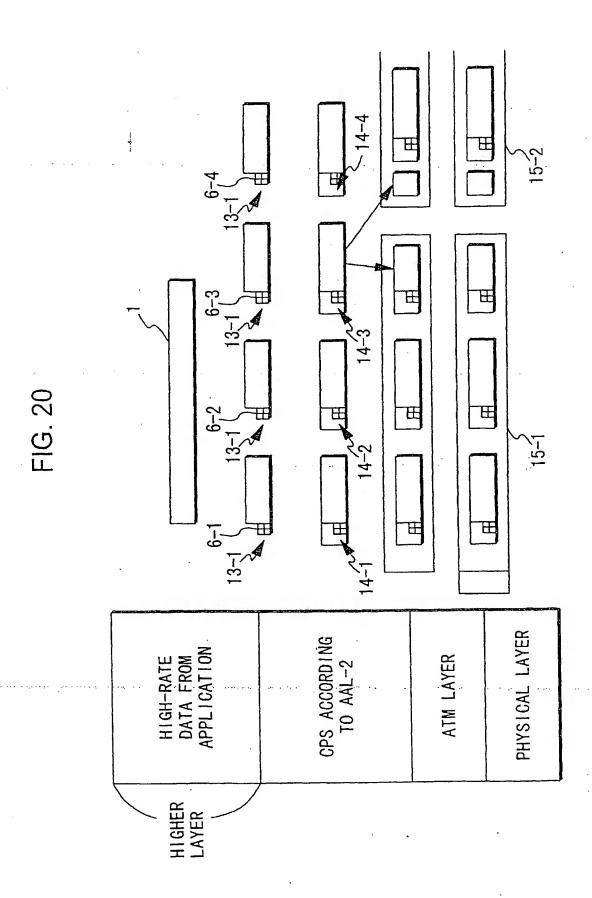
FIG. 16











# FIG.2

CODE FORMAT EXAMPLE OF THREE BIT UUI IN CPS PACKET HEADER.

000: SINGLE CPS PACKET ORIGINATED FROM SINGLE APPLICATION FRAME HAVING LENGTH EQUAL TO OR LESS THAN 45 OCTETS 001-011: CPS PACKET ORIGINATED FROM SEGMENT OF APPLICATION FRAME UU1 =010 HAVING LENGTH MORE THAN 45 OCTETS UU = -011 =001 MIDDLE END: 000= APPLICATION FRAME 010: 011: 011: 001: CPS ACCORDING

### INTERNATIONAL SEARCH REPORT

International application No.
PCT/JP97/04833

	SIFICATION OF SUBJECT MATTER C1 H04L12/56				
According to International Patent Classification (IPC) or to both national classification and IPC					
B. FIELD	S SEARCHED				
Minimum documentation searched (classification system followed by classification symbols) Int.Cl <sup>6</sup> H04L12/56, H04L12/28, H04L1/00-1/24					
Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched Annual Meeting of IEICE 1995-1997 Research Report of IEICE 1995-1997  Communication Society Convention of IEICE 1995-1997					
Electronic data base consulted during the international search (name of data base and, where practicable, search terms used)  JICST: "AAL" "ATM ADAPTATION LAYER"  INSPEC, COMPENDEX, WPI: "ATM" "AAL" "ATM ADAPTATION LAYER" "AAL2"  "TYPE 2"					
C. DOCU	MENTS CONSIDERED TO BE RELEVANT	·	<u> </u>		
Category*	Citation of document, with indication, where ap	propriate, of the relevant passages	Relevant to claim No.		
X Y	JP, 7-30543, A (NEC Corp.), January 31, 1995 (31. 01. 95	) (Family: none)	21-24, 26 25		
X Y	JP, 7-264192, A (Toshiba Con October 13, 1995 (13. 10. 95		21-24, 26 25		
Y	ITU-T Recommendation I.363 ( See : 4.AAL type 3	Helsinki, 1993)	25		
Y	JP, 6-52884, B2 (Nippon Telec and others), July 6, 1994 (06. 07. 94) (F	_	25		
A	1995 Fourth IEEE Internation Universal Personal Communica JAPAN) p.863-867 H. Nakamura carry very low bit-rate mobi 10.11.95)	tions Record (Tokyo, et al., "Using ATM to	1-20		
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Further documents are listed in the continuation of Box C. See patent family annex.					
"A" docume consider "E" earlier docume cited to special "O" documemans "P" documens	entegories of cited documents: ent defining the general state of the art which is not ired to be of particular relevance document but published on or after the international filing date ent which may throw doubts on priority claim(s) or which is establish the publication date of another citation or other reason (as specified) ent referring to an oral disclosure, use, exhibition or other ent published prior to the international filing date but later than ority date claimed	"T" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention document of particular relevance, the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone document of particular relevance, the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art document member of the same patent family			
	actual completion of the international search th 30, 1998 (30. 03. 98)	Date of mailing of the international search report April 14, 1998 (14. 04. 98)			
Name and mailing address of the ISA/ Japanese Patent Office		Authorized officer			
Facsimile No.		Telephone No.			

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### INTERNATIONAL SEARCH REPORT

International application No.
PCT/JP97/04833

C (Continua	ation). DOCUMENTS CONSIDERED TO BE RELEVANT		·
Category*	Citation of document, with indication, where appropriate, of the relevant	Relevant to claim No.	
<b>A</b>	B314 Masatomo Nakano et al., "Study on Sho Multiplex System in Mobile ATM Communication (in Japanese)", 1996 Communication Society Convention of IEICE (30.08.96)	Network	1-20
A	B-335 Hajime Tamura et al., "Layered-Cell Protcol Stack for Next-Generation Mobile Communication Network (in Japanese)", 1996 Annual Meeting of IEICE (11.03.96)		1-20
Α .	SB-12-7 Hiroki Yamada et al., "Study on Met Constructing Terminal Access CLAD realizing V Japanese)", 1996 Annual Meeting of IEICE (11	1-20	
A	SSE96-25 Takenori Okutani et al., "Study on ATM Celluating Method for Low-Speed Audio (in Japanese)" Research Report of IEICE (24.05.96)		1-20
P, A	B-5-208 Shuichi Tsuboya et al., "AAL type 2 Transmission System in Mobile ATM Communica Network (in Japanese)", 1997 Annual Meeting of (06.03.97)	tion [	1-20
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